
Easy to Win, Hard to Master: Playing Infinite Games Optimally

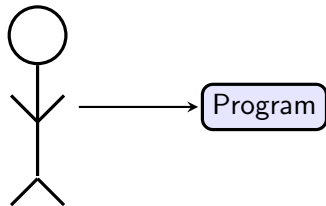
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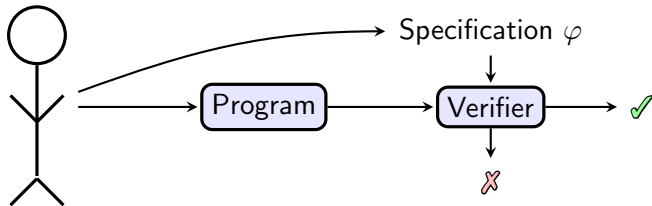
April 26th, 2017

Thesis Proposal Talk

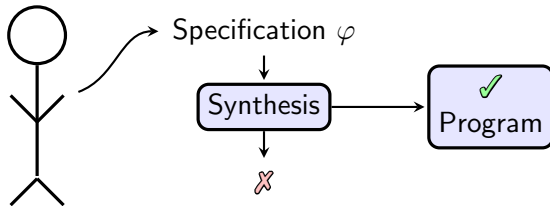
Programming



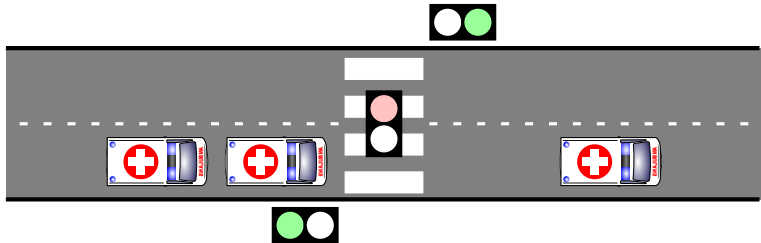
Program Verification



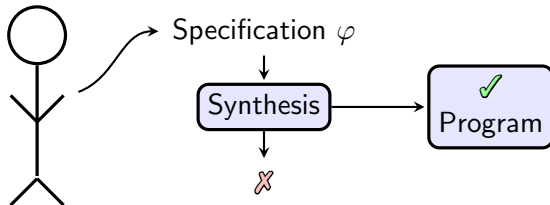
Program Synthesis



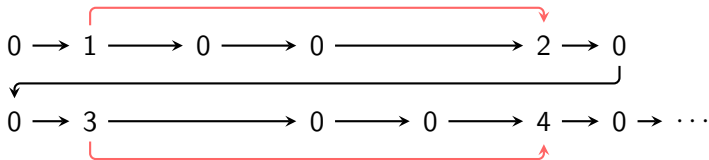
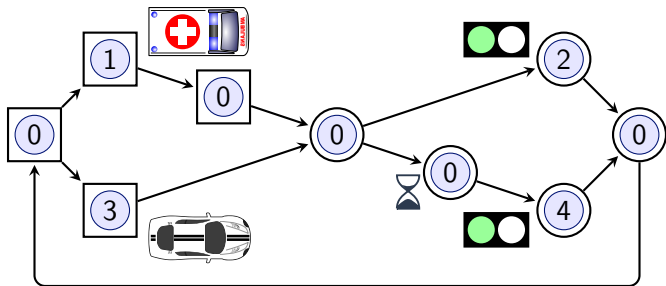
Example



Program Synthesis



Parity Games

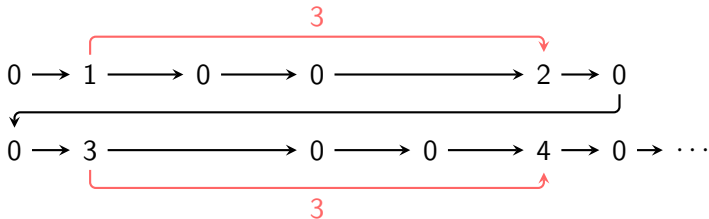
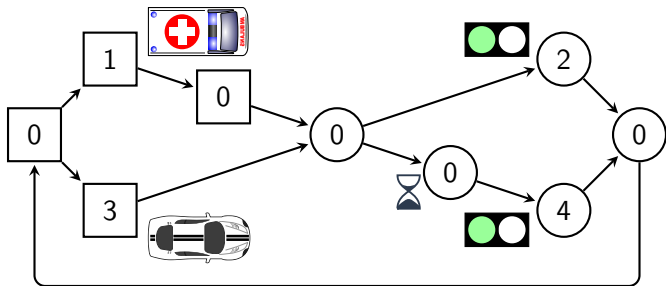


Example due to (Fijalkow and Chatterjee, Infinite-state games, 2013)

■ Deciding winner in $NP \cap co-NP$

■ Positional Strategies

Finitary Parity Games



Goal for Player 0: Bound response times

Decision Problem

Theorem (Chatterjee, Henzinger, Horn, 2009)

The following decision problem is in PTIME:

Input: Finitary parity game \mathcal{G}

Question: Does there exist a strategy σ with $\text{Cst}(\sigma) < \infty$?

Theorem (W., Zimmermann, 2016)

The following decision problem is PSPACE-complete:

Input: Finitary parity game \mathcal{G} , bound $b \in \mathbb{N}$

Question: Does there exist a strategy σ with $\text{Cst}(\sigma) \leq b$?

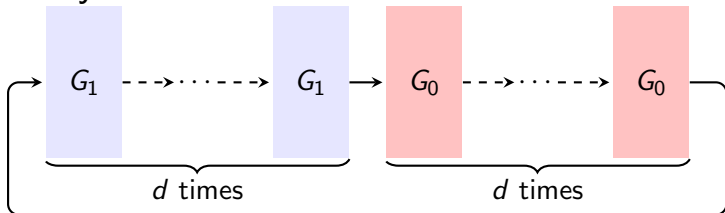
Memory Requirements (for Player 0)

Theorem (W., Zimmermann, 2016)

Optimal strategies for finitary parity games need exponential memory

Sufficiency: Corollary of proof of PSPACE-membership

Necessity:



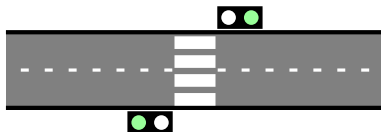
Player 0 needs to recall d positions with d possible values
 \Rightarrow Player 0 requires $\approx 2^d$ many memory states

Results so far

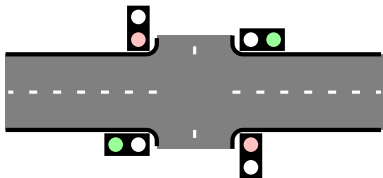
	Parity	Finitary Parity	
		Winning	Optimal
Complexity	$NP \cap co-NP$	P_{TIME}	P_{SPACE} -comp.
Strategy Size	1	1	Exp.

Outlook

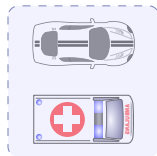
So far



Multi-Dimensional Games



Imperfect Information



Conclusion

Results so far: Forcing Player 0 to answer quickly in (finitary) parity games makes it harder

- to decide whether she can satisfy the bound
- for her to play the game

Guiding Question: What costs does playing games optimally incur

- in terms of computing a strategy?
- in terms of the complexity of strategies?